

## IN THE CLAIMS

Please amend the claims as follows:

1. – 8. (withdrawn).

9. (currently amended) A split-mask, masking countermeasure method for improving the resisting security attacks on resistance, to power analysis attacks, of a processing unit performing a defined cryptographic function using a key ~~to perform a defined cryptographic function~~, the method comprising the following steps:

obtaining the key and a random key mask value  $r$ ;

obtaining a set of  $n$  random input values  $m_{in1}, \dots m_{inN}$ ;

defining a masked function by masking the defined cryptographic function with the value  $m_{in1} \wedge \dots \wedge m_{inN}$ ;

masking the key with the random key mask value  $r$  to define the value  $mkey$ ;

obtaining a set of random split mask values  $m1, \dots mn-1$ ;

defining a split mask value  $mn$  to be  $r \wedge m_{in1} \wedge \dots \wedge m_{inN} \wedge m1 \wedge \dots \wedge mn-1$ ; and

using the values  $m1, \dots, mn$  and  $mkey$  to define input for the masked function.

10. (original) The method of claim 9 in which the encryption function is a table look-up.

11. (original) The method of claims 9 or 10 in which masking is a bitwise exclusive or operation carried out on binary values.

12. (currently amended) A split-mask, masking countermeasure method for improving the resistance, to power analysis attacks, resisting of a processing unit performing a cryptographic function ~~security attacks on a processing unit using a key to encrypt a plaintext value using a look up on a defined look-up table~~, the method comprising the following steps:

obtaining the key and a random key mask value r;

defining a value mkey by masking the key with the random key mask value r;

obtaining a set of n random input values  $m_{in1}, \dots, m_{inN}$ ;

defining a masked table by masking the defined look-up table with the value  $m_{in1} \wedge \dots \wedge m_{inN}$ ;

masking the key with the random value r to define the value mkey;

obtaining a set of split mask values comprising random values  $m1, \dots, mn-1$ ;

defining a split mask value  $mn$  to be  $r \wedge m_{in1} \wedge \dots \wedge m_{inN} \wedge m1 \wedge \dots \wedge mn-1$ ; and

masking the plaintext with the split mask values  $m1, \dots, mn$  and mkey to define input for the masked table, the masked table to be used in place of the defined look-up table in the cryptographic operation.

13. (original) The method of claim 12 in which masking is a bitwise exclusive or operation carried out on binary values.

14. – 29. (withdrawn).

30. (currently amended) A computing device program product for improving the resistance, to power analysis attacks, of a processing unit resisting security attacks on a processing unit using a key to perform a defined cryptographic function, the computing device program product comprising a computer usable storage medium having computer readable program code means embodied in said storage medium, and comprising

program code means for obtaining the key and a random key mask value r,

program code means for obtaining a set of n random input values  $m_{in1}, \dots, m_{inN}$ ,

program code means for defining a masked function by masking the defined cryptographic function with the value  $m_{in1} \wedge \dots \wedge m_{inN}$ ,

program code means for masking the key with the random key mask value r to define the value mkey,

program code means for obtaining a set of random split mask values m1, ... mn-1,

program code means for defining a split mask value mn to be  
 $r \wedge m_{in1} \wedge \dots \wedge m_{in n} \wedge m1 \wedge \dots \wedge mn-1$ , and

program code means for using the values m1,...,mn and mkey to define input for the masked function.

31. (original) The computing device program product of claim 30 in which the encryption function is a table look-up.
32. (original) The computing device program product of claims 30 and 31 in which masking is a bitwise exclusive or operation carried out on binary values.
33. (currently amended) A computing device program product for improving the resistance, to power analysis attacks, of a processing unit performing a cryptographic function resisting security attacks on a processing unit using a key to encrypt a plaintext value using a look up on a table, the computing device program product comprising a computer usable storage medium having computer readable program code means embodied in said storage medium, and comprising

program code means for obtaining the key and a random key mask value r,

program code means for obtaining a set of n random input values  $m_{in1}, \dots m_{in n}$ ,

program code means for defining a masked table by masking the defined look-up table with the value  $m_{in1} \wedge \dots \wedge m_{in n}$ ,

program code means for masking the key with the random key mask value r to define the value mkey,

program code means for obtaining a set of random split mask values  $m_1, \dots, m_{n-1}$ ,

program code means for defining a split mask value  $m_n$  to be

$r \wedge m_{in1} \wedge \dots \wedge m_{in} \wedge m_1 \wedge \dots \wedge m_{n-1}$ , and

program code means for masking the plaintext with the values  $m_1, \dots, m_n$  and mkey to define input for the masked table.

34. (original) The computing device program product of claim 33 in which masking is a bitwise exclusive or operation carried out on binary values.

35. – 58. (withdrawn)